

Fixing the Meaning of Object-Oriented Languages

Kim Bruce
Williams College
on leave at U. of California at Santa Cruz

What makes O-O Languages Special?

- Objects encapsulate state & methods
- Subtyping
- Inheritance
- Self-reference
 - Provides dynamic method invocation with more power!

2

Luca Cardelli:
"If there is no 'self' in a language then it is not object-oriented."

What's the Big Deal?

- Are objects more than records with function components?
- What provides real power?
- How can semantics and type theory help?
- Focus on class-based O-O languages like Smalltalk, Eiffel, & Java
 - Multi-method languages are quite different

4

Defining a Class



```
public class Squares {
    private FramedRect outer, inner;
    public Squares(Location upleft, int size,
        DrawingCanvas canvas){...}
    public void move(int dx, int dy) {
        outer.move(dx,dy);
        inner.move(dx,dy);
    }
    public void moveTo(int x, int y) {
        this.move(x - outer.getX(),
            y - outer.getY());
    }
}
```

5

Instance Variables

```
public class Squares {
    private FramedRect outer, inner;
    public Squares(Location upleft, int size,
        DrawingCanvas canvas){...}
    public void move(int dx, int dy) {
        outer.move(dx,dy);
        inner.move(dx,dy);
    }
    public void moveTo(int x, int y) {
        this.move(x - outer.getX(),
            y - outer.getY());
    }
}
```

6

Constructor

```
public class Squares {
    private FramedRect outer, inner;
    public Squares(Location upleft, int size,
        DrawingCanvas canvas){...}
    public void move(int dx, int dy) {
        outer.move(dx,dy);
        inner.move(dx,dy);
    }
    public void moveTo(int x, int y) {
        this.move(x - outer.getX(),
            y - outer.getY());
    }
}
```

7

Methods

```
public class Squares {
    private FramedRect outer, inner;
    public Squares(Location upleft, int size,
        DrawingCanvas canvas){...}
    public void move(int dx, int dy) {
        outer.move(dx,dy);
        inner.move(dx,dy);
    }
    public void moveTo(int x, int y) {
        this.move(x - outer.getX(),
            y - outer.getY());
    }
}
```

8

Creating & Using Objects

```
Squares fst = new Squares(corner,10,canvas);
Squares snd = new Squares(middle,40,canvas);
// objects are references
fst.moveTo(20,30);
snd = fst; // snd & fst refer to same object
snd.move(30,50);
```

9

Objects Are Fixed Points

First naive view of objects:

`[[new Squares (...)]] =`

```
μ this.({ outer = ..., // no mention of this
        inner = ... } ×
    { move = fun(dx,dy). this.outer.move()...,
      moveTo = fun(x,y). this.move(...) })
```

Defines mutually recursive methods.

10

Classes Are Generators

- Classes serve many roles:
 - Types
 - Generators of new objects
 - Extensible to form new generators

11

Subclass



```
public class OvalSquares extends Squares {
    private FramedOval center;
    public OvalSquares(Location upleft,
        int size, DrawingCanvas canvas) {
        super(upleft, size, canvas);
        center = new FramedOval(...);
    }
    public void move(int dx, int dy) {
        super.move(dx, dy); // old move
        center.move(dx, dy);
    }
} What happens to moveTo?
```

12

Classes Are Generators of Fixed Points

- Meaning of *this* is not bound in classes
 - Semantics of `moveTo` changes (indirectly) in `OvalSquares`
- class `Squares` = $\lambda \text{ this. SQ}(\text{this})$
- class `OvalSquares` = $\lambda \text{ this. OSQ}(\text{this})$
where `OSQ` extends `SQ`.
- Objects formed as *fixed points* of `SQ` and `OSQ`.

13

Objects From Subclasses

- `sq = new Squares(...);`
 - `sq = $\mu \text{ this. SQ}(\text{this})$`
- `osq = new OvalSquares(...);`
 - `osq = $\mu \text{ this'. OSQ}(\text{this}')$` // meaning of *this* changed!
 - where `super = SQ(this')` // uses new *this'* in body

14

Subtyping

- $T <: U$ iff any object of type `T` can *masquerade* as object of type `U`
- More formally, *subsumption* rule:
$$T <: U \ \& \ o : T \Rightarrow o : U$$
- Java Interfaces & extension

15

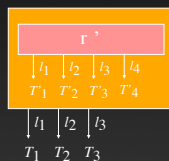
Subtyping Immutable Record Types

Records *without* field update: only operation is extracting field: ... `s.filling` ...

```
{bread: BreadTp; filling: CheeseTp; sauce: SauceTp}
<:
{ bread: BreadType; filling: FoodType }
iff
CheeseTp <: FoodType
```

16

Subtyping Immutable Record Types



$\{l_i : T'_i\}_{1 \leq i \leq n} <: \{l_i : T_i\}_{1 \leq i \leq k}$
iff
 $k \leq n$ and for all $1 \leq i \leq k$, $T'_i <: T_i$.

17

Subtyping Function Types

If $f : S \rightarrow T$ and $s : S$ then $f(s) : T$

When is $S' \rightarrow T' <: S \rightarrow T$?

If $f' : S' \rightarrow T'$ and $s : S$, need $f'(s) : T$.

18

Subtyping Function Types



$$S' \rightarrow T' <: S \rightarrow T$$

iff

$$S <: S' \text{ and } T' <: T.$$

Contravariant for parameter types.
Covariant for result types.

19

Subtyping Object Types

$$\text{ObjType } \{ m_i : T'_i \}_{1 \leq i \leq n} <: \text{ObjType } \{ m_i : T_i \}_{1 \leq i \leq k}$$

iff

$$k \leq n \text{ and for all } 1 \leq i \leq k, T'_i <: T_i,$$

assuming methods not updatable at run-time!

Method parameters can vary *contravariantly*,
 return types *covariantly*.

20

Restriction on Subclass Changes

- Java 1.4 doesn't allow any changes to types of methods in subclass.
- C++ & Java 5 allow covariant changes to return types.
- Suppose you don't care if subclass gives a subtype. Do you still need restrictions?
 - *In Smalltalk, subclass and subtype hierarchies sometimes reversed.*

21

```
class Example {
    :
    void m(...) {... this.n(s) ...}
    T n(S x) {...}
}

class SubExample extends Example {
    T' n(S' x) {...}
    void newMeth(...) {...}
}
```

What is relationship of new type of n to old if want type safety?

Restriction on Subclass Changes

- Method type in subclass must be subtype of method type in superclass for safety:
 - *Covariant* change allowed in return type
 - *Contravariant* change in parameter type

23

Semantics of Classes?

- Methods must be meaningful in all possible subclasses.

$[[\text{class } (i : I, m : M)]]$ =

$$\forall M' <: [[M]]. \forall IR' <: [[I^{ref}]].$$

$$[[i]] \times (\lambda(\text{this} : IR' \times (IR' \rightarrow M')). [[m]])$$

initial values of instance variables *methods*

24

Semantics of Objects

$[[\text{new Squares}(\dots)]]$ =

$\{ \text{outer} = \text{ref } \dots, \text{inner} = \text{ref } \dots \}$

\times

$\mu(\text{fm} : [[I^{\text{ref}}]]$

$\lambda(\text{inst} : [[I^{\text{ref}}]])$

$\{ \text{move} = \text{fun}(dx, dy). \text{inst}.\text{outer} \dots,$

$\text{moveTo} = \text{fun}(x, y). \langle \text{inst}, \text{fm} \rangle.\text{move}(\dots) \}$

Method suite can be shared between objects of same type. 25

Sending Messages

$[[\text{obj}]].\text{p}(\dots)$ = $\text{fm}(i).\text{p}(\dots)$

where $[[\text{obj}]]$ = $\langle i, \text{fm} \rangle$

In objects, methods fixed – parameterized by suite of instance variables, not *this*.

26

Summary

- *Fixed points* are key to understanding O-O languages.
- Classes are *extensible generators* of fixed points.
- *Subtyping* explains restrictions on subclasses
 - Even though subtyping distinct concept.

27

What is Type of "this"?

$[[\text{class}(i : I, m : M)]]$ =

$\forall M' \leq : [[M]]. \forall IR' \leq : [[I^{\text{ref}}]]$

$[[i]] \times \lambda(\text{this} : IR' \times (IR' \rightarrow M')). [[m]]$

How do we type function returning *this*?

28

Finding a Type for "this"

- Defining *this* in $\text{Class}(i : I, m : M)$
 - $\text{this} : IR' \times (IR' \rightarrow M')$ where $M' \leq : [[M]]$, $IR' \leq : [[I^{\text{ref}}]]$
- When return *this*, hide instance variables
 - $\text{this} : \exists X. X \times (X \rightarrow M') = \text{Obj}(M')$
 - but notice $\text{Obj}(M') \leq : \text{Obj}(M)$, thus $\text{this} : \text{Obj}(M)$
 - where $\text{Obj}(M)$ is type of objects of the class being defined.

29

Losing Information

- But now we have lost information.
- If C has method $\text{m}()$ with body `return this`, then $C.\text{m}()$.
- If D extends C, then inherited m still has type $C.\text{m}()$.
- But $\text{d.m}()$ really returns value of type D!

30

Introduce Type of "this"

- `ThisClass` is type of "this"
- In previous example, declare:
 - `m:() → ThisClass`
- Also examples with parameters of type `ThisClass` -- called *binary methods*

31

Binary Methods

```
public class Node {
    protected ThisClass next;
    public Node(ThisClass next) {
        this.next = next;
    }
    public void setNext(ThisClass newNext) {
        this.next = newNext;
    }
    public ThisClass getNext() {
        return next;
    }
}
```

32

Subclasses

```
public class DbleNode extends Node {
    protected ThisClass prev;
    public DbleNode(ThisClass next, ThisClass prev) {
        super(next);
        this.prev = prev;
    }
    public void setNext(ThisClass newNext) {
        super.setNext(newNext);
        newNext.setPrev(this);
    }
    public ThisClass setPrev(ThisClass newPrev) {
        this.prev = newPrev;
    }
    ...
}
```

33

But ...

There are type safety problems! Define:

```
public void breakIt(Node n1, Node n2) {
    n1.setNext(n2);
}
```

Let `n` be a `Node` and `dn` be a `DbleNode`:
`breakIt(dn, n);`

Blows up when `dn.setNext(n)` is executed!

Need to rule out this code using type system.

34

Fix Type-Safety

- Can only send binary message to object if know its exact type.
- If obj has type `@T`, then value of obj has type `T`, but not a subtype.

35

Using Exact Types

Rewrite `Node` and `DbleNode` to replace all occurrences of `ThisClass` by `@ThisClass`:

```
public void breakIt(@Node n1, @Node n2) {
    n1.setNext(n2);
}
```

More generally w/ type parameters could write:

```
public <N extends Node>
    void breakIt(@N n1, @N n2) {
        n1.setNext(n2);
    }
```

36

Meaning of ThisClass

- **ThisClass** is type of **this**.
- In semantics:
 - Node = μ **ThisClass**.Nd(**ThisClass**)
 - DbleNode = μ **ThisClass'**.DbNd(**ThisClass'**);
- Thus
 - **ThisClass** = Nd(**ThisClass**) = Node
 - **ThisClass'** = DbNd(**ThisClass'**) = DoubleNode

37

Type-Checking

- Wish to type-check modularly.
 - no repeated type-checking of inherited methods
- Type-check methods of class C under assumptions that hold in all extensions!

this : **ThisClass**

ThisClass extends C

Can prove soundness of type system.

38

Verification

- Verification can benefit from similar assumptions on **this** and **ThisClass**
- Alternative: Verify under closed-world assumption and copy inherited methods down.
 - *Compositional approach is better!*

39

There Is Much More ...

- Gets *much* more interesting when:
 - Allow type parameters (e.g., GJ)
 - Introduce ThisType as interface of **this**.
 - LOOJ extension of GJ (Java 5)
 - Allow mutually recursive ThisClasses in *class groups*.
 - Allows types and classes to be refined simultaneously in type-safe fashion.

40

Subject-Observer Example

```
SubjObsGrp = ClassGroup {
  class MySubject {
    void addObserver(MyObserver myObs) {...};
    void notifyObservers: proc(MyEvent myEvt) {...};
  }

  class MyObserver {
    void notify(MySubject mySubj, MyEvent myEvt) {...};
  }

  class MyEvent {...}
}
```

41

Extending Class Groups

```
ChoiceSubjObsGrp = ClassGroup extends SubjObsGrp {
  class extends MySubject {
    String getItem() {...};
  }

  class extends MyEvent {
    String getItem() {...};
  }

  class extends MyObserver {
    void notify(MySubject mySubj, MyEvent myEvt) {
      ... myEvt.getItem() ...;
    }
  }
}
```

42

Summary

- Fixed points play key role in OO languages.
- Classes are generators of fixed points:
 - $cl = \lambda \text{ this. } CL(\text{this})$
 - $obj = \mu \text{ this. } CL(\text{this})$
- Types similar:
 - $Class = \lambda \text{ ThisClass. } CLType(\text{ThisClass})$
 - $ObjTp = \mu \text{ ThisClass. } CLType(\text{ThisClass})$

43

Summary (continued)

- Type-checking must take into consideration future subclasses.
- Verification should also plan for subclasses.

44

Summary (continued)

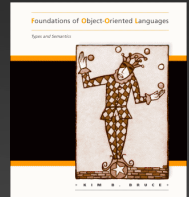
- Interdependent groups of objects common.
- Need to subclass simultaneously
 - Not supported in current languages
 - Mutually recursive ThisClasses provide opportunities for type-safe specialization.

45

References

- LOOJ: Weaving LOOM into Java, ECOOP 2004.
- Some Challenging Typing Issues in Object-Oriented Languages, ENTCS 82, #8, 2003.
- Foundations of Object-Oriented Languages: Types and Semantics, MIT Press, 2002.

<http://www.cs.williams.edu/~kim>



46

Questions?