

THEOREM

LET $F: \mathbb{B}^n \rightarrow \mathbb{B}$ BE A BOOLEAN FUNCTION.
 THEN ITS DUAL F^d IS GIVEN BY THE
 FORMULA

$$F^d(x_1, \dots, x_n) = \overline{F(\bar{x}_1, \dots, \bar{x}_n)} .$$

RMK: THE DUALITY PRINCIPLE FOLLOWS
 IMMEDIATELY FROM THIS RESULT SINCE THEN
 $F = G$ IFF $F^d = G^d$.

WE LEAVE THE PROOF OF THE ABOVE
 THEOREM AS AN EXERCISE (PROBLEM #27
 P. 708). HINT: USE THE RECURSIVE
 DEFINITION OF BOOLEAN EXPRESSION AND
 INDUCTION ON THE EXPRESSION WHICH
 REPRESENTS F . THE KEY STEP IS
 TO USE THE IDENTITIES $\overline{\bar{x}\bar{y}} = x+y$ AND
 $\overline{\bar{x}+\bar{y}} = xy$.

EX $F(x, y, z) = x\bar{y} + z\bar{x}$

$$\begin{aligned} \overline{F(\bar{x}, \bar{y}, \bar{z})} &= \overline{\bar{x}\bar{\bar{y}} + \bar{z}\bar{\bar{x}}} = (\overline{\bar{x}\bar{y}}) \cdot (\overline{\bar{z}\bar{x}}) \\ &= (\bar{\bar{x}} + \bar{y}) \cdot (\bar{\bar{z}} + \bar{x}) = (x + \bar{y})(\bar{z} + \bar{x}) \\ &= F^d(x, y, z) . \end{aligned}$$

10.2 REPRESENTING BOOLEAN FUNCTIONS

DEFN

A LITERAL is a BOOLEAN VARIABLE OR ITS COMPLEMENT: x OR \bar{x} . A MINTERM OF THE VARIABLES x_1, x_2, \dots, x_n IS A PRODUCT OF THE FORM

$$y_1 y_2 \dots y_n$$

WHERE EACH y_i IS EITHER x_i OR \bar{x}_i . I.e. A MINTERM IS A PRODUCT OF LITERALS

NOTE THAT A MINTERM IS TRUE FOR EXACTLY ONE COMBINATION OF VALUES OF ITS VARIABLES

EX $x\bar{y}\bar{z}$ IS A MINTERM OF x, y, z

x	y	z	\bar{y}	\bar{z}	$x\bar{y}\bar{z}$
0	0	0	1	1	0
0	0	1	1	0	0
0	1	0	0	1	0
0	1	1	0	0	0
1	0	0	1	1	1
1	0	1	1	0	0
1	1	0	0	1	0
1	1	1	0	0	0

OBSERVE THAT EVERY BOOLEAN FUNCTION CAN BE REPRESENTED AS A SUM OF MINITERMS OF ITS VARIABLES. THIS EXPRESSION IS CALLED THE SUM OF PRODUCTS REPRESENTATION OR DISJUNCTIVE NORMAL FORM (DNF) OF THE FUNCTION.

$$\begin{aligned}
 \text{EX. } F(x, y, z) &= \bar{x}(y + \bar{z}) \\
 &= \bar{x}y + \bar{x}\bar{z} \\
 &= \bar{x} \cdot y \cdot 1 + \bar{x} \cdot 1 \cdot \bar{z} \\
 &= \bar{x}y(z + \bar{z}) + \bar{x}(y + \bar{y})\bar{z} \\
 &= \bar{x}yz + \bar{x}y\bar{z} + \bar{x}y\bar{z} + \bar{x}\bar{y}\bar{z} \\
 &= \bar{x}yz + \bar{x}y\bar{z} + \bar{x}\bar{y}\bar{z}
 \end{aligned}$$

TO FIND THE DNF OF A FUNCTION, GIVEN ITS TRUTH TABLE, WRITE THE MINITERM CORRESPONDING TO EACH 1 IN THE OUTPUT COLUMN, THEN TAKE THE SUM OF THESE MINITERMS.

EX.	x	y	z	$G(x, y, z)$
	0	0	0	0
	0	0	1	1 $\longrightarrow \bar{x}\bar{y}z$
	0	1	0	0
	0	1	1	0
	1	0	0	0
	1	0	1	1 $\longrightarrow x\bar{y}z$
	1	1	0	1 $\longrightarrow xy\bar{z}$
	1	1	1	0

$$\therefore G(x, y, z) = \bar{x}\bar{y}z + x\bar{y}z + xy\bar{z}$$

THIS PROCEDURE SHOWS THAT ANY BOOLEAN FUNCTION CAN BE PLACED IN DNF, AND IN PARTICULAR THE MAPPING

$$\{\text{EXPRESSIONS}\} \rightarrow \{\text{FUNCTIONS}\}$$

IS ONTO. NOTE ALSO THAT THE DNF OF A FUNCTION IS UNIQUE UP TO ORDER OF TERMS.

LIKewise ANY BOOLEAN FUNCTION CAN BE REPRESENTED BY A PRODUCT OF SOME EXPRESSIONS CALLED THE CONJUNCTIVE NORMAL FORM (CNF).

THIS FORM CAN BE OBTAINED BY TAKING DOALS TWICE AND USING $F^{dd} = F$.

EX. $F(x, y) = xy + \bar{x}\bar{y}$

$$\begin{aligned} (xy + \bar{x}\bar{y})^d &= (x+y)(\bar{x}+\bar{y}) \\ &= x\bar{x} + \bar{x}y + x\bar{y} + y\bar{y} \\ &= 0 + \bar{x}y + x\bar{y} + 0 \\ &= \bar{x}y + x\bar{y} \end{aligned}$$

$$\begin{aligned} (xy + \bar{x}\bar{y})^{dd} &= (\bar{x}y + x\bar{y})^d \\ &= (\bar{x}+y)(x+\bar{y}) \end{aligned}$$

$$\therefore F(x, y) = (\bar{x}+y)(x+\bar{y})$$

SINCE EVERY BOOLEAN FUNCTION CAN BE REPRESENTED USING THE OPERATORS \cdot , $+$, $-$ WE SAY THAT THE SET

$$\{ \cdot, +, - \}$$

IS FUNCTIONALLY COMPLETE.

ACTUALLY THIS IS NOT THE SMALLEST SUCH SET. USING THE IDENTITY $x + y = \overline{\overline{x} \cdot \overline{y}}$ WE CAN ELIMINATE ALL OCCURRENCES OF $+$, SHOWING THAT $\{ \cdot, - \}$ IS ALSO FUNCTIONALLY COMPLETE.

LIKEWISE $x \cdot y = \overline{\overline{x} + \overline{y}}$ SHOWS THAT $\{ +, - \}$ IS FUNCTIONALLY COMPLETE AS WELL.

ITS ALSO TRUE THAT THERE ARE COMPLETE SETS CONSISTING OF JUST ONE OPERATOR.

DEEN

x	y	NAND	NOR
		$x \downarrow y$	$x \downarrow y$
0	0	1	1
0	1	1	0
1	0	1	0
1	1	0	0

OBSERVE THAT $\{|\}\}$ IS FUNCTIONALLY COMPLETE SINCE

$$\bar{x} = x | x \quad \& \quad x y = (x | y) | (x | y)$$

AND THE FACT THAT $\{\cdot, -\}$ IS FUNCTIONALLY COMPLETE.

LIKEWISE $\{\downarrow\}$ IS COMPLETE SINCE

$$\bar{x} = x \downarrow x \quad \& \quad x + y = (x \downarrow y) \downarrow (x \downarrow y)$$

AND THE FACT THAT $\{+, -\}$ IS COMPLETE.

EXERCISE: PROVE ALL OF THE ABOVE IDENTITIES.

EXERCISE:

FIND THE DNF OF $x | y$ AND $x \downarrow y$, AND SHOW THAT

$$(x \downarrow y)^{\text{df}} = x | y$$